

Please fill in your Student Number and Name.

Student Number : _____

Name:

Student Number:

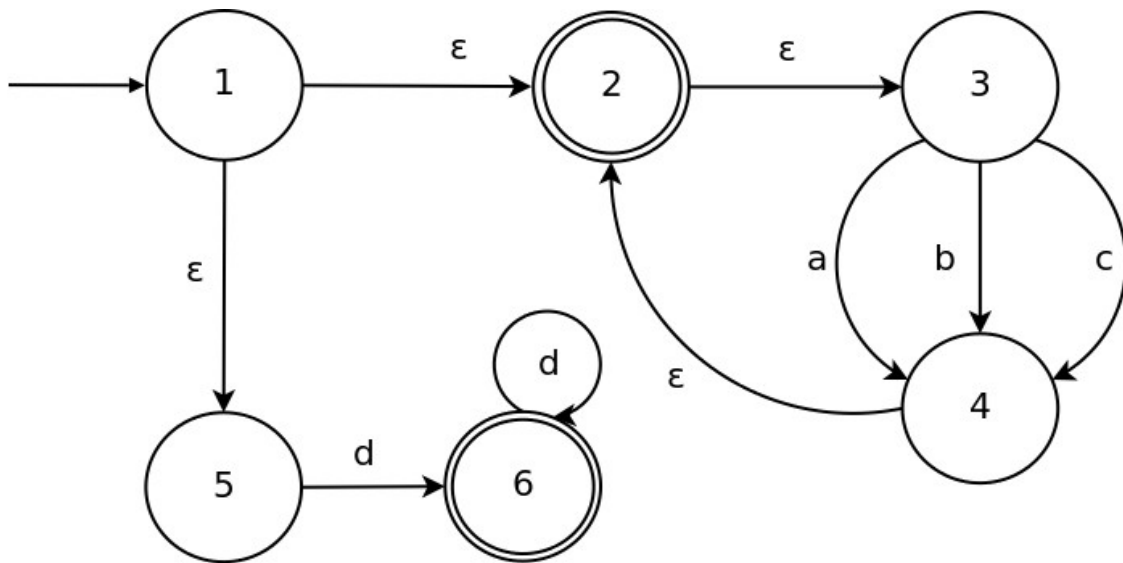
University of Cape Town ~ Department of Computer Science
Computer Science 3003S ~ 2009
Test 1

Question	Max	Mark	Internal	External
1	12			
2	10			
3	8			
4	15			
TOTAL	40			

Marks : 40
Time : 45 minutes
Instructions:

- a) Answer all questions.
- b) Write your answers in the space provided.
- c) Show all calculations where applicable.

c) Convert the Non-Deterministic Finite Automaton (NFA) below to a Deterministic Finite Automaton (DFA). Remember to indicate the start and final state(s). [7]

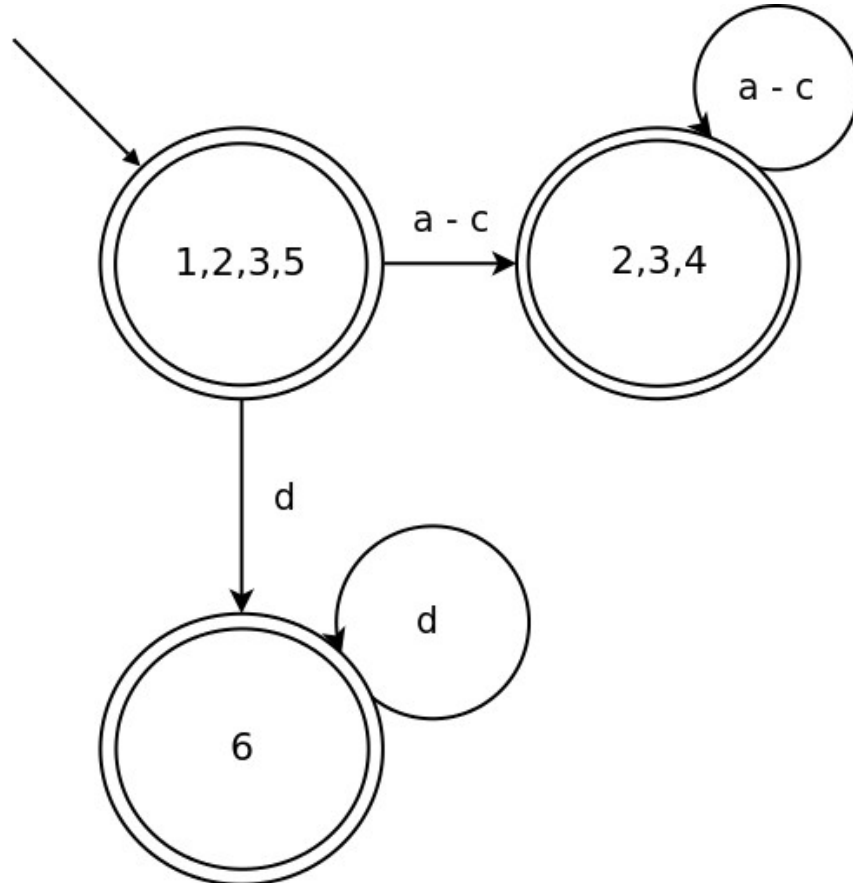


Solution DFA shown below

½ Mark for correct state (labels may differ), maximum 1½ marks

½ Mark for showing that state is final, maximum 1½ marks

1 Mark per correctly labelled transition (if student shows separate edges for a, b and c then only 1 mark is given for all of them), maximum 4 marks



Question 2 – Parsing [10]

a) Consider the following grammar and complete the table below by calculating its nullable, FIRST and FOLLOW sets (\$ should be treated as a terminal): [5]

- S → E\$
- E → aFc
- E → aF
- F → b
- F →

	nullable	FIRST	FOLLOW
S			
E			
F			

Solution table shown below

1 Mark for correct nullable column

1 Mark for correct FIRST column

1 Mark for correct FOLLOW entry for E

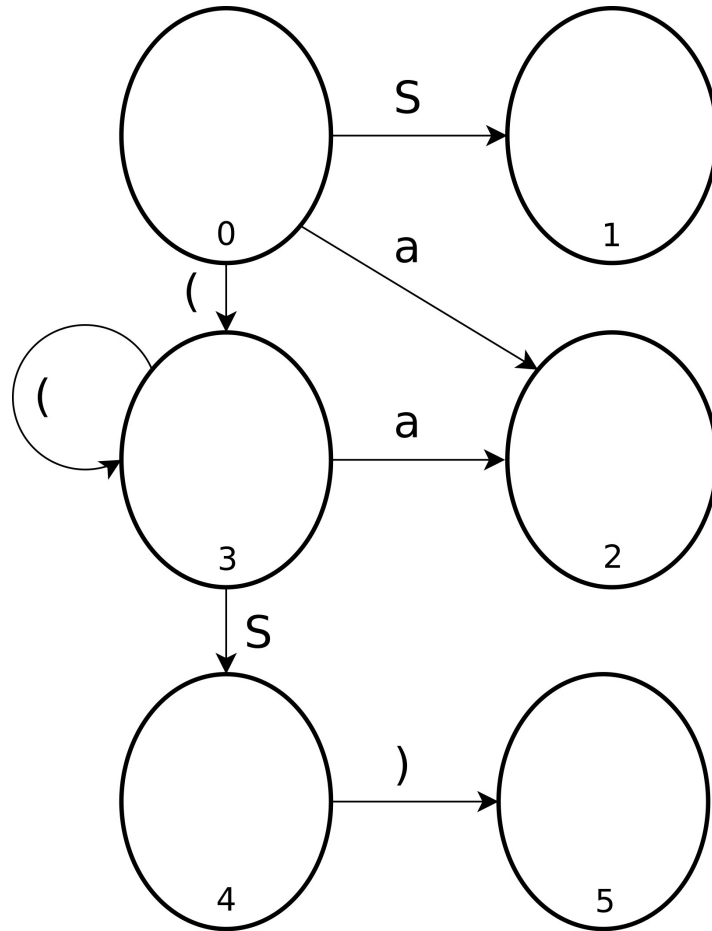
2 Marks for correct FOLLOW entry for F (ignore comas)

	<i>nullable</i>	<i>FIRST</i>	<i>FOLLOW</i>
<i>S</i>	<i>FALSE</i>	<i>a</i>	
<i>E</i>	<i>FALSE</i>	<i>a</i>	<i>\$</i>
<i>F</i>	<i>TRUE</i>	<i>b</i>	<i>c, \$</i>

b) Consider the following grammar and complete the LR(0) automaton below:

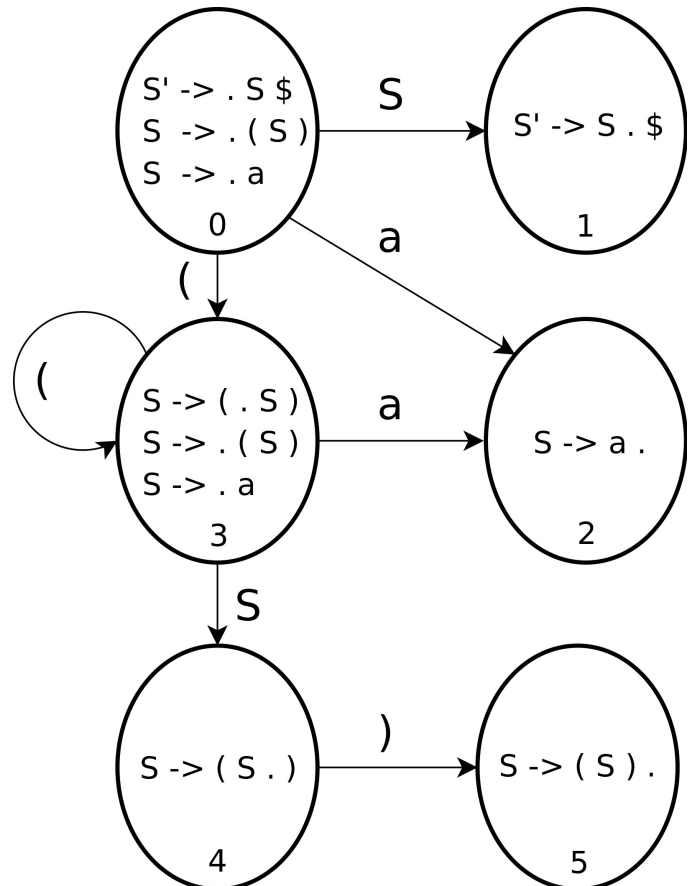
[5]

- $S' \rightarrow S\$$
- $S \rightarrow (S)$
- $S \rightarrow a$



Solution automaton on the right

½ Mark for each correct LR(0) item (production with a dot), max 10 items => 5 marks



Question 3 – Error Recovery [8]

There are two main classes of error recovery techniques:

- Local Error Recovery
- Global Error Recovery

a) What is a difference between these two classes? [2]

*Global Error Recovery adjusts the input before the point where the error was detected [1]
Local Error Recovery adjusts the input stream at the point of error (or immediately after) [1]*

b) What adjustments can be made to the input stream to enable parsing to continue? [3]

Tokens can be deleted [1], inserted [1] or replaced/substituted [1]

Burke-Fisher is an example of an error recovery technique.

c) Given a token window size of 5, how many alterations can be made to the input stream to enable parsing to continue? [1]

One [1]

d) For which edge case(s) does this technique fail (assuming there are no semantic actions that affect lexical analysis) and why? [2]

Burke-Fisher works by single token manipulation of the input stream [1], thus it fails if there is no single token adjustment that will allow parsing to continue [1].

i.e. 1 mark for stating single token manipulation and 1 mark for linking it to edge case

Question 4 – Semantic Analysis [15]

Semantic analysis follows syntactic analysis and is the process of determining context-sensitive information and thereafter generating code. Answer the following questions related to semantic analysis.

- a) The starting point in semantic analysis is usually an Abstract Syntax Tree. What is the difference between a Concrete Syntax Tree and an Abstract Syntax Tree? [2]

Concrete syntax tree contains all tokens from a grammar as derived from parsing. [1]

Abstract syntax trees contain equivalent semantic information in concise tree structures. [1]

- b) Symbols tables are used to perform context-sensitive checking. Name 2 types of information and 2 distinct attributes that may be stored in a symbol table. [2]

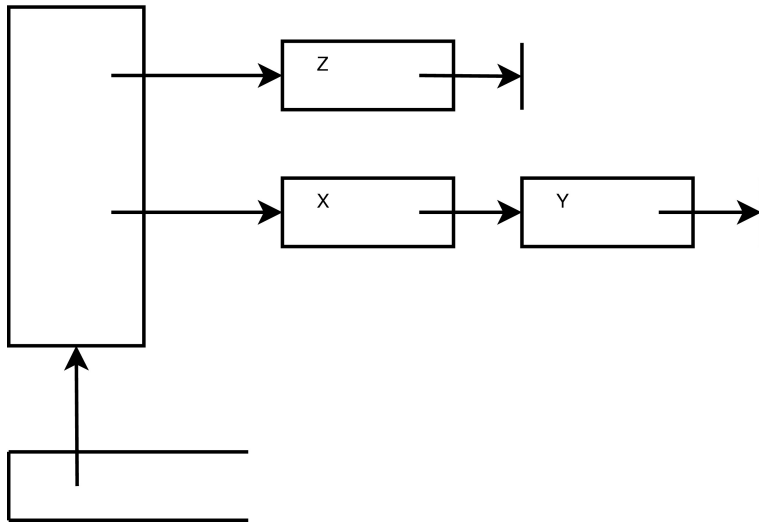
information: variable definitions, type definitions, subprogram names, etc. [1/2 x 2]

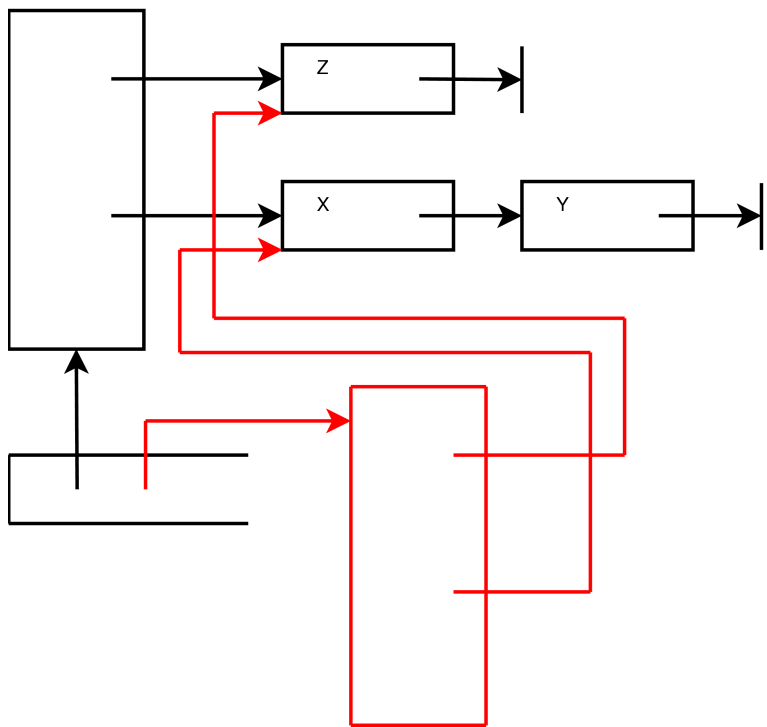
types: names, type, size, lexical level, etc. [1/2 x 2]

- c) Scope must be taken into account in a symbol table. How does static scope differ from dynamic scope? [2]

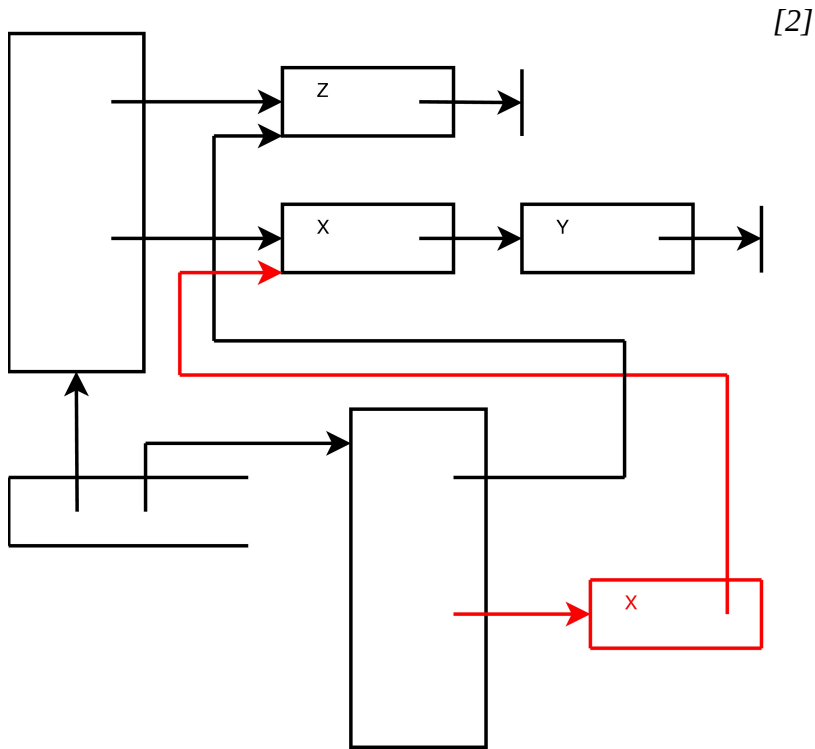
static scope is determined by lexical nesting of subprograms [1] – dynamic scope is determined by call sequence [1]

- d) Show diagrammatically how the following **functional hash table** implementation of a symbol table is modified (A) when a new scope is opened and (B) when a new definition of X is added in that scope. [4]





[2]



- e) Briefly explain 2 techniques to resolve non-local references at runtime in statically-scoped languages. [2]

follow static links to get to lexically-nested parents' activation records. [1]

use a display of links to subprograms at different lexical levels of nesting [1]

- f) The following DDD assembler source code listing corresponds to the unoptimised C function provided. In the context of subprogram invocation mechanics, circle and indicate which memory address or register use corresponds to: (A) a local variable; (B) the frame pointer; and (C) the return value. [3]

```

push    %ebp
mov     %esp,%ebp
sub     $0x10,%esp
mov     0x8(%ebp),%eax
add     $0x1,%eax
mov     %eax,-0x4(%ebp)
mov     -0x4(%ebp),%eax
leave
ret

```

```

int testsub ( int x )
{
    int a = x+1;

```

```
    return a;  
}
```

A – any -0x4(ebp) [1]
B – any ebp [1]
C – eax in last mov statement [1]